The *abc* **Conjecture of the Derived Logarithmic Functions of Euler's Function and Its Computer Verification**

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Abstract

Regarding Euler's (totient) function, for an arbitrary number $n > 1$, there exists a *k* that possesses the characteristic where $\varphi^k(n) = 1$. In this case, if *k* is expressed as $L(n)$ for *n*, then *L* possesses the characteristic of being perfectly logarithmic. For this *L*, we (Yamashita, Miyata) have provided the following *L* version *abc* conjecture.

Conjecture: When *a*, *b*, and *c* are relatively prime, numbers are natural, and $a + b = c$, then

 $\max\{L(a), L(b), L(c)\} \leq 2 \cdot L(rad (abc))$

is feasible.

This paper describes the properties of *L* and presents verification that this conjecture is correct up to $10⁹$ using a computer experiment. We also note that the *abc* conjecture recently considered solved by Prof. Mochizuki at Kyoto University is different from the conjecture presented here.

Introduction

Considering $\varphi^k(x) = \varphi(\varphi^{k-1}(x))$ $(k > 1)$ as $\varphi^1(x) = \varphi(x)$ with respect to Euler's function φ , when $x > 1$ then $\varphi(x) \leq x$. Therefore, there always exists a minimum *k* such that $\varphi^{k}(x) = 1$ for all $x > 1$. Heretofore, in regard to the properties of this *k*, Pilali ([1],[2]), Shapiro ([3],[9]), Murányi ([4]), et al have shown that *k* possesses (imperfect) logarithmic characteristics. Since then, a great deal of research on this has been conducted. Currently, it is known that by modifying this *k* (hereinafter, this *k* shall be indicated as $L(x)$), that the same becomes perfectly logarithmic.^{*1}

In this paper, we describe the properties and the extensions of the logarithmic function $L(x)$ derived of Euler's function and note that the *abc* conjecture pertaining to $L(x)$ we provide holds even under appropriate conditions other than primitive *φ*-triple, and also cite ours proof of this conjecture.

1. Various Properties of $L(x)$

1.1 Perfect logarithms of $L(x)$ **and the evaluation thereof**

Definition. 1. (Yamashita, [5]) *L is defined for the natural number n as follows and is called* a derived logarithmic function *of Euler's function*.

$$
L(n) = \begin{cases} 0 & (n = 1) \\ L(\varphi(n)) & (n: odd number > 1) \\ L(\varphi(n)) + 1 & (n: even number). \end{cases}
$$

At this time,

Proposition. 2. *L is perfectly logarithmic for any natural number x, y, i.e.*, $L(xy) = L(x) + L(y)$.

Therefore, the following simple evaluation can be obtained for *L*.

Proposition. 3. *If* $L(x) = n$, *then* $2^n < x < 3^n$ Then, immediately from there:

Corollary. 4. (E1) If $x \leq 2^n$ then $L(x) \leq n$. $(E2)$ *If* $x \geq 3^n$ *then* $L(x) \geq n$.

Corollary. 5. Let $x = 2^t \cdot x_0$ (x_0 : *odd*). If $L(x) = n$, then $x \leq 2^t \cdot 3^{n-t}.$

Corollary. 6. Let $x = 2^t \cdot x_0$ (x_0 : *odd*). If $x > 2^t \cdot 3^{n-t}$, then

 $L(x) > n$.

etc. can be obtained, and the following evaluation formula can also be obtained.

Proposition. 7.

(E3) $\log_3 2 \left(\min \left(L(x), L(y) \right) + 1 \right) \le L(x + y)$ \leq log₂ 3 (max $(L(x), L(y)) + 1$) $(E4) L(x - y) \le \log_2 3 \max (L(x), L(y))$ Remark: $log_2 3=1.58496250...$, $log_3 2 = 0.63092975...$

As for this *L*, we have also obtained the following theorem as an extension form of Euler's function *φ*.

Theorem. 8. (*Miyata*–*Yamashita*, [11], [12]) *Let* **P** *be a set of prime numbers and* $P \rightarrow N$ (*natural numbers*) *be a function such that* $1 \le f(p) < p \in P$. *If*

$$
\varphi_f(x) = x \prod_{i=1}^r \frac{f(p_i)}{p_i}, \ x = p_1^{e_1} p_2^{e_2} ... p_r^{e_r}
$$

and

$$
L_{\varphi f}(1) = 0
$$

$$
L_{\varphi f}(x) = L(\varphi_{\varphi f}(x)) + \#\{p \in f^{-1}(1) : p \mid x\}.
$$

then

$$
L_{\varphi f}(xy) = L_{\varphi f}(x) + L_{\varphi f}(y).
$$

The φ_f in the above theorem is a formal generalization of Euler's function by *f*. Also, according to the symbol of this theorem, $L(x) = L_a(x)$.

1.2 Extensibility of $L(x)$

L defined on the natural numbers can naturally be extended on $\mathbb{Z}\setminus\{0\}$ via *L*(-1) = 0, *L*(-*x*) = *L*(*x*). For *L*(0), if we define, for example, *L*(0) = ∞ , it

can also be is defined on **Z**. Therefore, if we define *L x* $\frac{1}{y}$ = *L*(*x*) – *L*(*y*) for *x* $\frac{\partial}{\partial y} \in \mathbf{Q}^{\times} = \mathbf{Q} \setminus \{0\}$, then we have a natural extension to **Q**. In other words,

the following holds:

Proposition. 9. *The L in Definition* 1 *can be naturally expanded on rational*

numbers **Q** *and the properties of Proposition* 2 *are also inherited.*

Can this *L* (here is where we part ways with the world of Euler's function *φ*) be expanded to a number $O[\sqrt{-1}]$ which is obtained by adding $\sqrt{-1}$ to real numbers **R** and **Q**, or complex number **C**, while maintaining the properties of Proposition 2?

Let us calculate by assuming the properties of Proposition 2. If we do some calculations with irrational numbers then,

$$
L(\sqrt{2}) = L(2^{1/2}) = \frac{1}{2}L(2) = \frac{1}{2}
$$

$$
L(\frac{1}{\sqrt{2}}) = L(2^{-1/2}) = -\frac{1}{2}L(2) = -\frac{1}{2}
$$

$$
L(2^{\sqrt{2}}) = \sqrt{2}L(2) = \sqrt{2}
$$

$$
L(2^{\pi}) = \pi L(2) = \pi
$$

When observing this situation, in order for *L* to be welldefined even on **R**, the range must be at least **R**.

In addition, let us continue to observe **C** as well. If

$$
\omega = \cos\left(\frac{2\pi}{n}\right) + \sqrt{-1}\sin\left(\frac{2\pi}{n}\right) \quad (n \in \mathbb{N})
$$

then from

$$
nL(\omega) = L(\omega^n) = L(1) = 0
$$

we obtain

$$
L(\omega)=0.
$$

In addition, if we let $\zeta = \cos \alpha + \sqrt{-1} \sin \alpha$, and then take *β* as $\alpha \beta = 2\pi$, then

$$
\beta L(\zeta) = L(\zeta^{\beta}) = L(1) = 0 \quad \text{if } L(\zeta) = 0
$$

Then, it will be $L(w) = 0$ for the point *w* on the unit circle of a complex plane, and the arbitrary *z* of C has the form $z = |z|w$. Therefore, we can obtain

$$
L(z) = L(|z|w) = L(|z|) + L(w) = L(|z|).
$$

However, there remain issues as to whether *L* can continue or extend in a well-defined manner from **Q** to **R** and **R** to **C** (including handling of transcendental numbers).

2. *abc* **Conjecture for the Derived Logarithmic Function** *L*

2.1 On the *abc* **conjecture for** *L*

Regarding this *L*, we have provided an *abc* conjecture (*L* version *abc* conjecture) pertaining to this derived logarithmic function *L* of Euler's function.

Conjecture. (Yamashita–Miyata [14])

Let a, b, c be coprime. If $a + b = c$ *, then* $\max\{L(a), L(b), L(c)\}$ < 2 · *L* (rad (*abc*)).

Regarding this conjecture, we confirmed the correctness up to $c < 2^{30}$ by computer verification (Miyata-Yamashita [16]), and by touching lightly on the proof of the polynomial version *abc* conjecture by Stothers ([8]). The results we obtained were as follows.

Theorem. 10. (Yamashita–Miyata [14]) *Let a, b, c be coprime. If* $a + b = c$, *then*

 $\max\{L(a), L(b), L(c)\}$ < 2 · *L*(rad (*abc*)).

The condition of Theorem 10 where $\varphi(a) + \varphi(b) = \varphi(c)$ is feasible, (a, b, c) *c*), is called **primitive** *φ***–triple** (Miyata–Yamashita [17]). Yamashita–Miyata have argued regarding the feasibility status of primitive *φ*–triple, and it is predicted to exist infinitely many times, and it is also known that the probability of existence of primitive *φ*–triple differs greatly due to the even/odd of *c* (Yamashita–Miyata [17]).

2.2 Cases other than primitive *φ***–triple**

In Theorem 10 we asserted that our conjecture is correct in the case of primitive *φ*–triple (Yamashita–Miyata [11]). However, what about cases other than primitive *φ*–triple?

Let *p* and *q* be coprime, and assume

$$
\frac{q}{p} = \frac{\varphi(a) + \varphi(b)}{\varphi(c)}
$$

If so, then the following theorem holds.

Theorem. 11. Let a, b, c be coprime. If $a + b = c > 2$, then under the following

condition (*)

(*) max
$$
(L(p), L(q)) \leq (2 - \log_2 3) L(\text{rad}(abc))
$$

\nmax $\{L(a), L(b), L(c)\} < 2 \cdot L(\text{rad}(abc)).$

we obtain

Proof. When simultaneously both
$$
a + b = c
$$
 and $p\varphi(a) + p\varphi(b) = q\varphi(c)$, then we obtain

$$
ac\left(q\frac{\varphi(c)}{c} - p\frac{\varphi(a)}{a}\right) = bc\left(p\frac{\varphi(b)}{b} - q\frac{\varphi(c)}{c}\right)
$$

Then if we assume

$$
q\frac{\varphi(c)}{c} - p\frac{\varphi(a)}{a} = 0
$$

then

$$
aq\varphi(c) = cp\varphi(a) = (a + b)\varphi(a)
$$

On the other hand, $q\varphi(c) = p\varphi(a) + p\varphi(b)$ results via

$$
(\# 1) \quad a\varphi(b) = b\varphi(a)
$$

However it must be $a|\varphi(a)$ because $(a, b) = 1$ and it must be $a = 1$. Meanwhile, if $a = 1$, then it is $\varphi(b) = b$ via (# 1), therefore $b = 1$, resulting in a contradiction in $2 < c = a + b = 1 + 1 = 2$. Therefore,

$$
q\frac{\varphi(c)}{c} - p\frac{\varphi(a)}{a} \neq 0.
$$

From which follows:

$$
\frac{a}{b} = \frac{p\frac{\varphi(b)}{b} - q\frac{\varphi(c)}{c}}{q\frac{\varphi(c)}{c} - p\frac{\varphi(a)}{a}}
$$

$$
= \frac{\text{rad}(abc)\left(p\frac{\varphi(b)}{b} - q\frac{\varphi(c)}{c}\right)}{\text{rad}(abc)\left(q\frac{\varphi(c)}{c} - p\frac{\varphi(a)}{a}\right)}
$$

$$
= \frac{p\left(\text{rad } (abc) \frac{\varphi(b)}{b}\right) - q\left(\text{rad } (abc) \frac{\varphi(c)}{c}\right)}{q\left(\text{rad } (abc) q \frac{\varphi(c)}{c}\right) - p\left(\text{rad } (abc) \frac{\varphi(a)}{a}\right)}
$$

Then, if we note the fact that with $k = a$, b, c then

$$
\operatorname{rad} \left(abc \right) \frac{\varphi(k)}{k} \in \mathbf{N}
$$

(hereinafter, rad (*abc*) will be denoted as rad*), then

$$
a | p \left(\text{rad}^* \frac{\varphi(b)}{b} \right) - q \left(\text{rad}^* \frac{\varphi(c)}{c} \right).
$$

Therefore,

$$
L(a) \le L\left(p\left(\text{rad}^*\frac{\varphi(b)}{b}\right) - q\left(\text{rad}^*\frac{\varphi(c)}{c}\right)\right).
$$

If the domain of *L* is expanded **Q** and we note that

$$
L\left(\frac{\varphi(k)}{k}\right) = \begin{cases} -1 & (k: \text{even})\\ 0 & (k: \text{odd}) \end{cases}
$$

regarding $k = a, b, c$, we can then use Proposition 7 (E4), which results in the following right side of the above equation:

$$
L\left(p\left(\text{rad}*\frac{\varphi(b)}{b}\right) - q\left(\text{rad}*\frac{\varphi(c)}{c}\right)\right).
$$

Simply, if we denote as rad* $\frac{\varphi(k)}{k} = C(k)$ (*k* = *a*, *b*, *c*) then

$$
L(a) \le \log_2 3 \cdot \max(L(pC(b)), L(qC(c)))
$$

\n
$$
\le \log_2 3\left(\left(\frac{2}{\log_2 3} - 1\right)L(\text{rad*}) + \max(L(C(b)), L(C(c)))\right)
$$

\n
$$
\le \log_2 3\left(\left(\frac{2}{\log_2 3} - 1\right)L(\text{rad*}) + L(\text{rad*})\right)
$$

\n
$$
= 2L(\text{rad*}) = 2L(\text{rad } (abc))
$$

In the case of primitive φ -triple, since $L(p) = L(q) = L(1) = 0$, then the conditions of Theorem 11 are satisfied and it can then be obtained as a corollary.

Corollary. 12. (Theorem 10) *If a* primitive *φ*–triple *then* $\max\{L(a), L(b), L(c)\}$ < 2 $\cdot L(\text{rad}(abc))$ /

3. Computer Verification of the Conjecture

3.1 The difficulty of computer verification for $c \le N = 10^{10}$

For our conjecture, computer experiments have confirmed that the conjecture is true for $c < 2^{30}$ (Mivata–Yamashita [15]), but with $c \ge 2^{30}$ and above it is difficult to verify using a typical PC environment.

Generally, the problem of finding $\varphi(x)$ for *x* is called an RSA problem, and if $\varphi(x)$ is easily obtained, the RSA public key encryption problem terminates, hence this is a very challenging problem.

Since it is necessary to repeatedly calculate $\varphi(x)$ to calculate $L(x)$, finding $L(x)$ involves more difficulty than the RSA problem.

On the other hand, if $L(x)$ is found for all x where $O(N \log \log N)$, it is known that time complexity $O(N \log \log N)$ can be used [15]. With that method, $L(x)$ can be obtained with $O(\log \log N)$ per each case.

However, this method requires a storage area for $O(N)$ which amounts to 4 GB of memory for $N = 10^9$.

In order to execute $N = 10^{10}$ in the same way (since the integers to be handled exceed 32 bits, it would mean using a 64-bit integer type), 80 GB of memory is required, which is impossible to execute on a typical PC.

As follows, verification was performed at $c \leq N = 10^{10}$ for (1, *b*, *c*). The verification results are shown in Table 1, and $q(1, b, c)$ in the Table is called a quality of (1, *b*, *c*), expressed as

$$
q(1,b,c) = \frac{L(c)}{L(\text{rad}(bc))}.
$$

The verification environment was as follows:

- \bullet PC: Acer Veriton X4620G
- OS: Windows 8.1 Pro
- CPU: Intel Core *i*5-3340 CPU (3.10GHz)
- \bullet RAM: 12.0GB
- Language: Java 9.0.1 (64-bit) Java (TM) SE Develoment Kit 9.0.1 (64-bit)
- Software: Eclipse

Execution time: 36 minutes 54 seconds

3.2 Computer verification for $c < 10^{10}$ **regarding** $(1, b, c)$

3.2.1 memoization

In the verification of $(1, b, c)$, $L(x)$ and $L(\text{rad}(x))$ were calculated in advance for $x \le 10^8$ using the Miyata–Yamashita method ([15]). As necessary, for reference, memoization was implemented. The storage area required for this is about 800 MB.

To obtain $L(x)$ for $x > 10^8$, the function was first factorized into prime numbers by trial division to obtain $\varphi(x)$, and then the memo could be referenced with values less than 10^8 . When $x > 10^8$, we sought to the greatest extent possible not to evaluate $L(x)$.

3.2.2 Finding the maximum prime of *c*

In the verification algorithm, for $S = 10^6$, the calculation was performed by dividing 1010 into segments of size *S*.

For example, for the *k*–th segment, calculation is performed for $c = kS + k$ 1, $kS + 2, \ldots, (k + 1)$ *S*, and then, using the segmented sieve algorithm, the largest prime factor of each *c* was sought.

If $S \le \sqrt{N}$, $N(i.e., O(S \log N)$ per each one), $O(S \log N)$ is sufficient for the calculation amount. Also, the storage area was $O(S)$ (in reality, 16*S* bytes $= 16$ MB required).

Let *p* be the largest prime factor of and $c = xp$. If $p < 10^8$, $L(c)$ and $L(\text{rad}(c))$ can be computed at high speed.

The reason being, if $p > 100$ then $x < 10^8$, it is therefore sufficient to merely reference the memo for both $L(x)$ and $L(p)$ because $L(c) = L(x) +$ *L*(*p*).

On the other hand, if $p < 100$, c can be factorized into prime numbers at high speed because it is rendered with the product of small prime numbers less than 100.

3.3 $(1, b, c)$ –triple determination

Definition. 13. Let $1 + b = c$. $(1, b, c)$ *that satisfy*

$$
\frac{L(c)}{L(\text{rad}(abc))} \ge 1.25
$$

is called (1, *b*, *c*)–**triple**.

In this verification, we will judge whether each c is

$$
\frac{L(c)}{L(\text{rad}\,(p-1))+L(\text{rad}\,(c))} \ge 1.25
$$

Determination is conducted as follows, with *p* being the maximum prime factor of *c*, and *q* being the maximum prime factor of $c - 1$.

3.3.1 Case $p \ge 10^8$

If $c = p$, then $L(p) / (L(\text{rad}(c-1)) + L(p)) < 1$. If $c = xp$, then $1 \le x \le 100$. Therefore,

$$
\frac{L(c)}{L(\text{rad}(c-1)) + L(\text{rad}(c))}
$$
\n
$$
\leq \frac{L(x) + L(p)}{2 + L(p)}
$$
\n
$$
= 1 + \frac{-2 + L(x)}{2 + L(p)}
$$
\n
$$
\leq 1 + \frac{-2 + \log_2 10^2}{2 + \log_3 10^8} < 1.248.
$$

Therefore, in this case, $(1, c - 1, c)$ does not become a triple.

3.3.2 Case *p* **< 108 and** *q* **< 108**

In this case, $L(c)$, $L(\text{rad}(c))$, $L(c-1)$, $L(\text{rad}(c-1))$ can be calculated at high speed. So we actually calculate as

$$
\frac{L(c)}{L(\text{rad}(c-1))+L(\text{rad}(c))}
$$

and then we simply need to investigate whether it is 1.25 or higher or not.

3.3.3 Case *p* **< 108 and** *q* **\$ 108** Since $L(q) \ge \log_3 10^8$, then we first seek out

$$
\frac{L(c)}{L(\text{rad}(c))+1+\log_3 10^8}
$$

If this is not 1.25 or higher, then we can determine that is not a triple.

If not, then we calculate $L(c-1)$ and $L(\text{rad}(c-1))$ while factorizing into prime numbers, then determine whether it is

$$
\frac{L(c)}{L(\text{rad}(c-1)) + L(\text{rad}(c))} \ge 1.25
$$

or not.

By the way, the number of cases where it was necessary to calculate $L(c -$ 1) and $L(\text{rad}(c-1))$ by prime factorization was only several hundred times out of $c \le 10^{10}$. Of those, those that were 1.25 or higher more were 0 times.

3.3.4 The reason for a threshold of 1.25

There are three reasons why we used 1.25 as the sieving threshold for the verification algorithm.

Reason1: The lower the threshold, the higher the number of corresponding triples. And, for this study, we were not interested in small triples.

Reason2: Our conjecture was

$$
\frac{\max\{L(c-1), L(c)\}}{L(\text{rad}(c-1))+L(\text{rad}(c))} < 2
$$

but the enumeration is

$$
\frac{L(c)}{L(\operatorname{rad}\left(c-1\right))+L(\operatorname{rad}\left(c\right))} \ge 1.25.
$$

If there is a *c* such that

$$
\frac{\max\{L(c-1), L(c)\}}{L(\operatorname{rad}(c-1))+L(\operatorname{rad}(c))} \ge 2,
$$

then it will be

$$
\frac{L(c)}{L(\text{rad}(c-1)) + L(\text{rad}(c))} \ge 2 \log_3 2 > 1.26
$$

 which means that it will always be included in this enumeration. Therefore, setting the threshold to 1.25 makes it possible to verify that there is no counterexample.

Reason3: As a practical reason, we tried memoization for 108 or less as we wanted to be able to be execute this verification using a personal computer of ordinary specifications. $(10⁹$ is impossible to do without a slightly high-performance personal computer.)

4. Summary and Future Issues

In this study, we examined the domain extensibility of $L(x)$, further improved results using primitive *φ*–triple, and showed our conjecture is correct for non-primitive *φ*–triples as long as certain conditions were met.

In terms of verifying our conjecture, we focused on $(1, b, c)$ and verified that our conjecture is correct for $C \leq 10^{10}$.

In terms of future issues, we still need a proof for our conjecture's feasibility, but we also need to verify our conjecture. For the time being, we will further increase N until $c \leq N$. However, of note are the following:

- \bullet In the case of $(1, b, c)$, we will increase the evaluation accuracy of the inequality in Section 4.5.1 and lower the sieve threshold to below 1.248.
- We will optimally apply the inequality condition of Theorem 11 to the verification algorithm.

This paper is a partial addition to [19].

Notes

*1. Yamashita showed that *k* according to a different definition from theirs was completely logarithmic in his high school days ([5]). After that, in 1977, during correspondence with Professor Saburo Uchiyama (Tsukuba University) (Yamashita-Uchiyama, Uchiyama-Yamashita [6],[7]) he learned for the first time of Pilali ([1],[2]), Shapiro ([3]), and Murányi's work ([4]). However, at this timing, facts in a perfect logarithmic form were not known in academic circles. It was not perfectly logarithmic in the first edition of Shapiro's textbook in 1983 ([9]). The first time it become known that it was perfectly logarithmic in academic circles was in the note made by Prasad, et al. ([10]).

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Calculation Results

b	C	q(1, b, c)
$5^4 \cdot 19 \cdot 15541$	$2^{24} \cdot 11$	1.35
$3^{13} \cdot 1277$	$2^3 \cdot 19^2 \cdot 89^3$	1.35
$2^3 \cdot 7^5 \cdot 13^2 \cdot 109$	$3^3 \cdot 11^3 \cdot 41^3$	1.35
$2^5 \cdot 3^2$	17 ²	1.333
$2^5 \cdot 3 \cdot 5^2$	7 ⁴	1.333
$3^2 \cdot 5 \cdot 7 \cdot 13$	2^{12}	1.333
$3^4 \cdot 79$	$2^8 \cdot 5^2$	1.333
$5 \cdot 11^3$	$2^9 \cdot 13$	1.333
$2^6 \cdot 3^2 \cdot 5 \cdot 29$	17 ⁴	1.333
$2^5 \cdot 3 \cdot 5 \cdot 7 \cdot 29^2$	41 ⁴	1.333
$5^3 \cdot 7^4 \cdot 11$	$2^{13} \cdot 13 \cdot 31$	1.333
$7^4 \cdot 2399$	$2^{10} \cdot 3^2 \cdot 5^4$	1.333
$2^5 \cdot 3^3 \cdot 7 \cdot 13 \cdot 307$	17 ⁶	1.333
$3^7 \cdot 53 \cdot 131^2$	$2^{20} \cdot 7 \cdot 271$	1.333
$3^9 \cdot 5^4 \cdot 709$	$2^6 \cdot 53 \cdot 137^3$	1.333
$2^6 \cdot 3^{10} \cdot 331$	$17^5 \cdot 881$	1.318
$7^2 \cdot 71^2 \cdot 223$	$2^{15} \cdot 41^2$	1.316
$2^{14} \cdot 8111$	$3^5 \cdot 5^7 \cdot 7$	1.313
$7^3 \cdot 487$	$2 \cdot 17^{4}$	1.308
$7^2 \cdot 13^2 \cdot 186391$	$2^{26} \cdot 23$	1.304
$19^3 \cdot 23^2 \cdot 1613$	$2^{19} \cdot 3 \cdot 61^2$	1.304
$2^{12} \cdot 5^3$	$3^5 \cdot 7^2 \cdot 43$	1.3
$31^3 \cdot 79^2$	$2^{16} \cdot 2837$	1.3
$3^7 \cdot 11 \cdot 19^2 \cdot 31$	$2^{18} \cdot 13 \cdot 79$	1.3
$3^2 \cdot 7^3 \cdot 19^4$	$2^{13} \cdot 49109$	1.3
$7^4 \cdot 13 \cdot 23^2 \cdot 59$	$2^4 \cdot 3^6 \cdot 17^4$	1.3
$5 \cdot 139^{3}$	$2^7 \cdot 3 \cdot 11^2 \cdot 17^2$	1.294
$2^7 \cdot 3 \cdot 5^2 \cdot 7^4$	4801 ²	1.294
$2^3 \cdot 3^7 \cdot 5^4 \cdot 7$	$13^2 \cdot 673^2$	1.294
$2 \cdot 3 \cdot 281^{3}$	$7^5 \cdot 89^2$	1.294

The *abc* Conjecture of the Derived Logarithmic Functions of Euler's Function and Its Computer Verification

b	C	q(1,b,c)
$3^5 \cdot 643$	$2 \cdot 5^7$	1.25
$5 \cdot 7^4 \cdot 19$	$2^8 \cdot 3^4 \cdot 11$	1.25
$3 \cdot 5^2 \cdot 11 \cdot 31 \cdot 41$	2^{20}	1.25
$5 \cdot 29 \cdot 47^{3}$	$2^9 \cdot 3^5 \cdot 11^2$	1.25
$2^4 \cdot 5^2 \cdot 7^2 \cdot 13^2 \cdot 29$	$3^8 \cdot 11^4$	1.25
$5^9 \cdot 163$	$2^4 \cdot 3^3 \cdot 23 \cdot 179^2$	1.25
$3^2 \cdot 7 \cdot 11 \cdot 31 \cdot 151 \cdot 331$	2^{30}	1.25
$3^8 \cdot 13^2 \cdot 2311$	$2^{18} \cdot 5^2 \cdot 17 \cdot 23$	1.25
$2^2 \cdot 5^4 \cdot 17^3 \cdot 211$	$3^3 \cdot 7^3 \cdot 23^4$	1.25

The *abc* Conjecture of the Derived Logarithmic Functions of Euler's Function and Its Computer Verification